

General Conversion Scheme of Card-based Protocols for Two-colored Cards to Updown Cards

— the number of cards for
computing an arbitrary function

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tentative edition



Takumi Sakurai

- Takumi Sakurai is a student of the Graduate School of Informatics in Nagoya University, Japan.
- His research is about card-based cryptography.

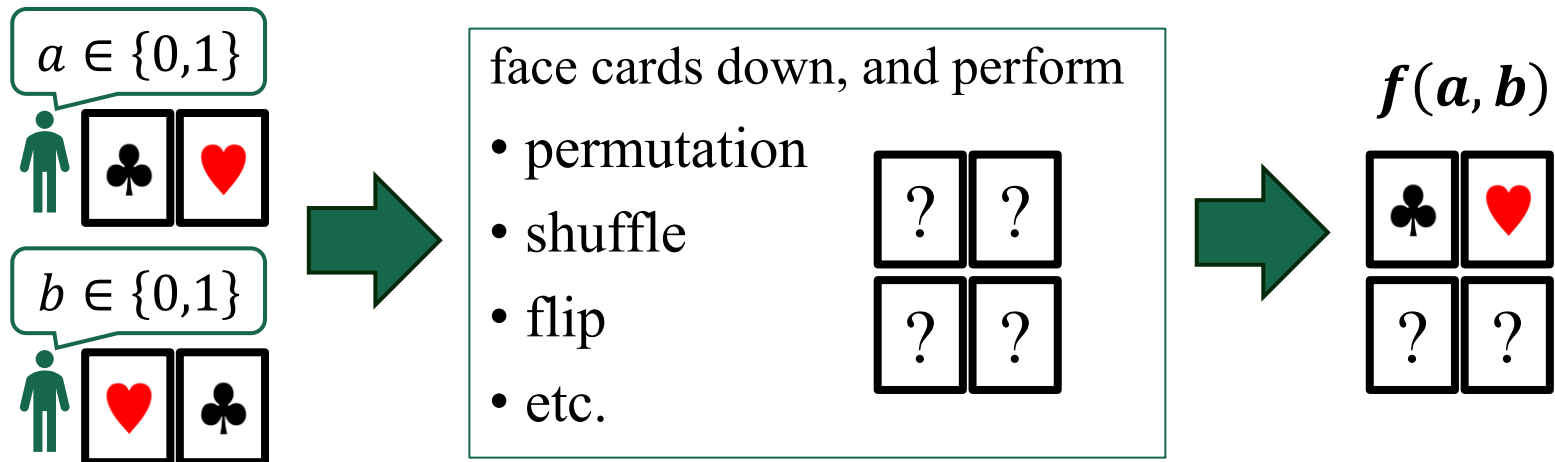
Summary

In card-based cryptography, we will ...

- discuss the sufficient condition of two-colored cards protocols that can be converted to updown cards.
- clarify the number of updown cards that compute an arbitrary Boolean function by the conversion.
- This study is a compilation of many known results, rather than a novel proposal based on a new idea.

Card-based cryptography

- a technique for secure multi-party computation using physical cards



- without specialized knowledge or equipment
- lectures of security, zero-knowledge proof for puzzles [Gradwohl 07][Shinagawa 22]

Example: AND protocol [Mizuki 09]

- two-colored cards

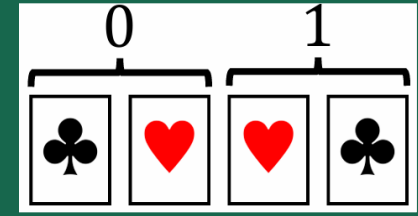
- printed with either “♣” or “♥” on front and “?” on back
front:   back:  

- Cards with the same symbol are not distinguished.
- placed on a table. We can see only the upper side of cards.
- A bit is encoded as a *commitment*, a pair of two cards.

$$\begin{array}{|c|} \hline \clubsuit \\ \hline \end{array} \begin{array}{|c|} \hline \heartsuit \\ \hline \end{array} = 0, \quad \begin{array}{|c|} \hline \heartsuit \\ \hline \end{array} \begin{array}{|c|} \hline \clubsuit \\ \hline \end{array} = 1$$

- The negation bit is obtained by swapping cards in a commitment.

Example: AND protocol [Mizuki 09]

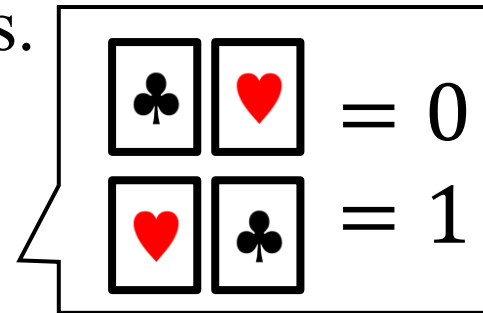
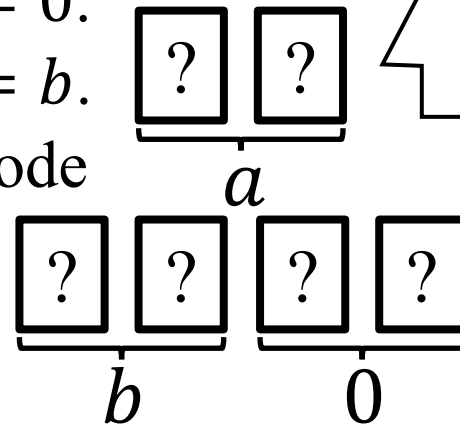


- Alice has $a \in \{0,1\}$. Bob has $b \in \{0,1\}$.
- They obtain only the result of $a \wedge b$.
- Alice and Bob want to compute secure AND.

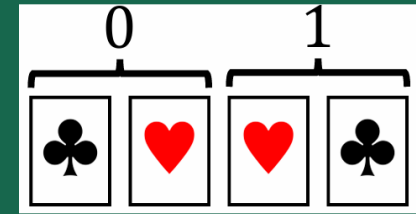
protocol's procedure

1. Put the input cards and additional cards.

- Keep the cards face down.
- When $a = 0$, $a \wedge b = 0 \wedge b = 0$.
- When $a = 1$, $a \wedge b = 1 \wedge b = b$.
- The cards under the heart encode the result of the computation.

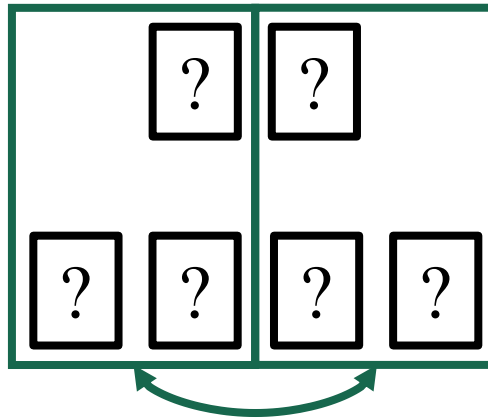


Example: AND protocol [Mizuki 09]



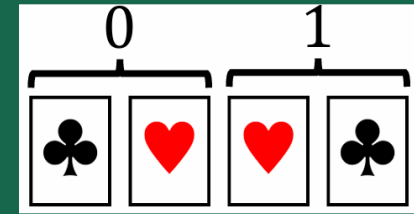
2. Shuffle the left and right sides of the cards.

- We can use card sleeves or clips.

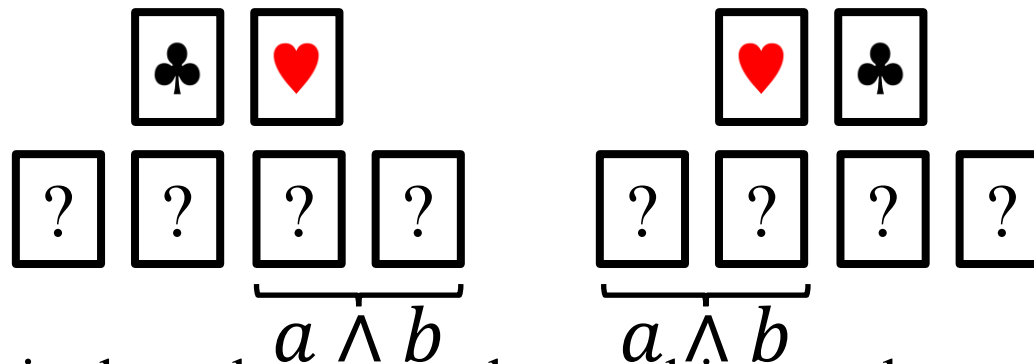


- Participants don't know whether the cards are swapped.
- Probability of swapping the cards is $1/2$.

Example: AND protocol [Mizuki 09]



3. Open the top two cards, and determine the position of the result.
- The result is the commitment under the heart.
 - By the shuffle, the bit a is not identified.



- To obtain the value, open the resulting cards.
- We can keep the resulting cards face down, and use them for the subsequent computation.

Two-colored protocols

- 6 cards AND protocol [Mizuki 09]
- 4 cards XOR protocol [Mizuki 09]

These basic protocols construct the below protocols

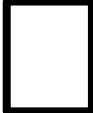
[Nishida 14]

- $2n + 6$ cards
arbitrary n -variable Boolean function
- $2n + 2$ cards
arbitrary n -variable symmetric Boolean function

Updown cards

- Updown cards

- printed with a single symbol of an arrow

- front:  back: 

- A bit is encoded by the direction of the arrow.

$$\begin{array}{|c|} \hline \downarrow \\ \hline \end{array} = 0, \quad \begin{array}{|c|} \hline \uparrow \\ \hline \end{array} = 1$$

- The negation bit is obtained by rotating the card.
 - Updown protocols require fewer cards than two-colored.
 - Investigations for updown cards is not enough.
 - There is no updown protocols computing an arbitrary Boolean function.

Main idea

- Given an arbitrary updown protocol, it is clearly possible to construct two-colored protocol. [Shinagawa 23]
 - Retain the expressions of bits.



- Two-colored protocol that preserves commitments can be converted to updown protocol.



Conversion scheme

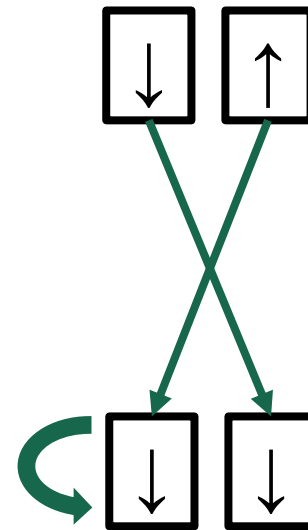
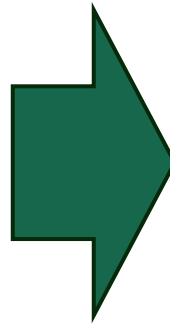
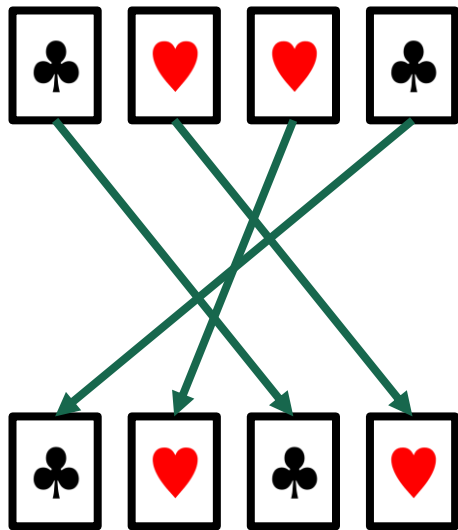
The two-colored protocol using only the following operations can be converted to updown protocol with half the number of cards.

- permutation not destroying the commitment
- shuffle not destroying the commitment
- flip that turns a card face up or face down
- NOT operation

These operation can be simulated by updown cards.

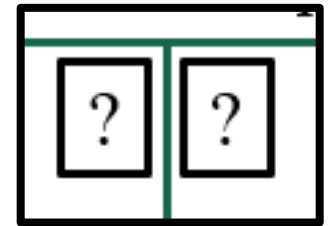
Example of conversion

- with two-colored cards
 - permutation not destroying the commitment
- with updown cards
 - permutation and NOT operations



Logical operations

- Two-colored protocols computing an arbitrary Boolean function destroy the commitment.
 - AND and XOR protocols have the operations separating two cards of a bit.



- With updown cards, AND and XOR protocols were proposed with half the number of cards.
 - 3 cards AND protocol [Mizuki 14][Shinagawa 16]
 - 2 cards XOR protocol [Mizuki 14]

Conversion scheme

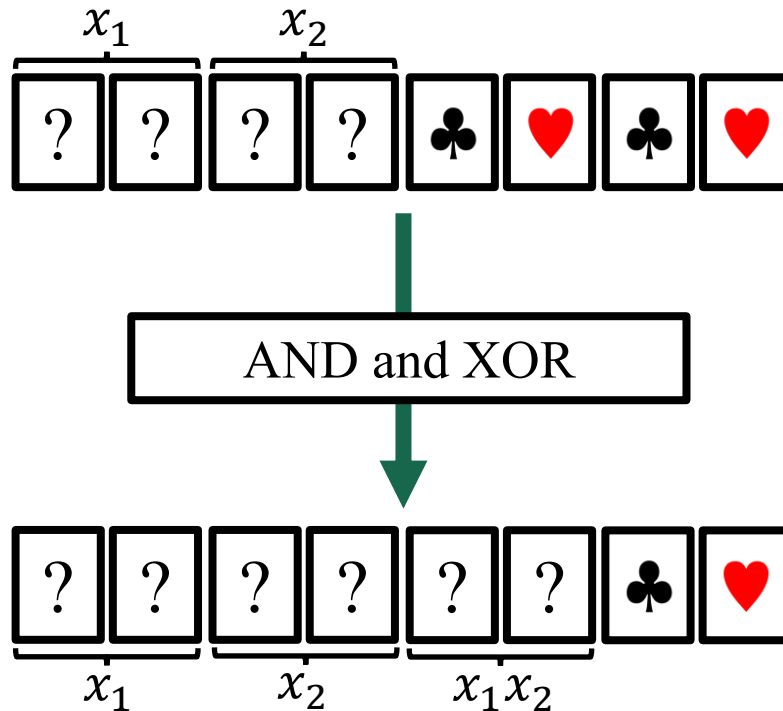
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- 4 cards XOR protocol

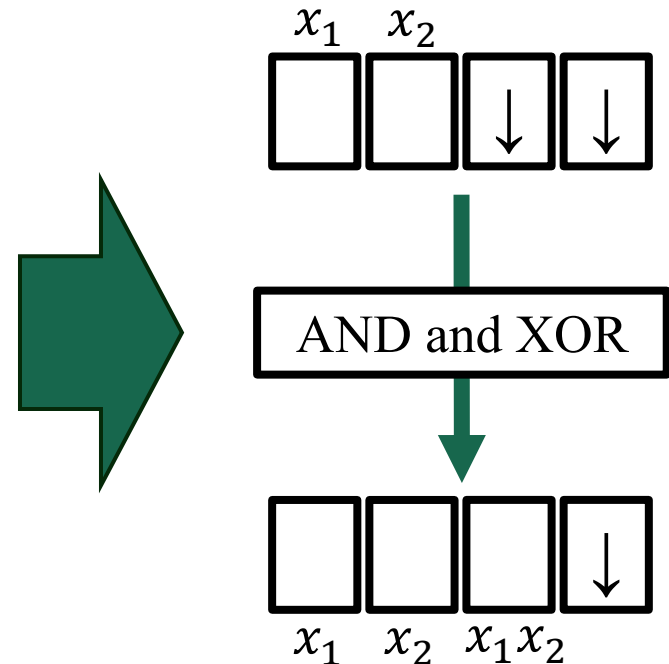
Example of applications

- with two-colored cards arbitrary n -variable Boolean function

$f(x_1, x_2, \dots, x_n)$ [Nishida 14]



- with updown cards



Arbitrary Boolean function

These basic protocols construct the below protocols

[Nishida 14]

- $2n + 6$ cards
arbitrary n -variable Boolean function
- $2n + 2$ cards
arbitrary n -variable symmetric Boolean function

Convert two-colored protocols to updown protocols.

- $n + 3$ cards
arbitrary n -variable Boolean function
- $n + 1$ cards
arbitrary n -variable symmetric Boolean function

Conclusion and future work

- Conclusion
 - conversion scheme for the two-colored protocol to the updown protocol
 - Boolean function with $n + 3$ updown cards
symmetric function with $n + 1$ updown cards
- Future work
 - about other protocols destroying the commitment
 - The minimum number of cards is still unresolved.
 - Use the relationship between updown and two-colored cards...

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Thanks for your attention!