



#### On Compressing Time-Evolving Networks

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#### Presenter's Bio

- Sudhindra Gopal Krishna is a Ph.D. Candidate in the School of Computer Science at the University of Oklahoma.
- Sudhindra graduated with a Master's Degree in Computer Science in Spring 2017.
- Current research interests:
  - Algorithms, Data Compression, Graphs, Network Science.





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#### Introduction

- A Graph (aka Network) is a structure represented as a set of nodes (or vertices) connected by a set of links (or edges).
- A Graph is formally represented as G = (V, E).
  - V represents a non-empty set of vertices, and
  - *E* represents a set of edges.
- Each edge in the graph is represented as a tuple (u, v).





## Time-Evolving Graphs

- Time-evolving graphs can be stored as an extension to the static graphs.
  - By adding the third and the fourth dimensions, which are the start time  $(t_i)$ , and the end time  $(t_j)$ .
  - Formal notation of a time-evolving graph is  $G_i = (V_i, E_i, \tau)$ ,
    - where *i* is a time-frame, and
    - $\tau$  is a time interval  $(t_i, t_j)$ .
  - So, a link in a time-evolving graph is a 4tuple  $(u, v, t_i, t_j)$ .
- This is often seen in social networks.





# Massive Graphs

- The graphs that do not fit into the memory for analysis are referred to as massive graphs.
- In reality, these graphs are very sparse.
- Where **Sparsity** is a measure of how many links exist out of all possible links in the graph.



# Storage Solution

- One way to address the issue of storing massive graph is to compress the data.
- There are two ways to compress a graph
  - Graph (Matrix) based compression, and
  - Node (Row-by-Row) based compression.
- All the algorithms in this paper are based on row-by-row compression.







# Background – Node Based Compression

- Time-evolving graphs can also be stored as a series of static graphs called snapshots.
- Row-by-Row compression for static graphs,
  - Compressed Sparse Row (CSR) [1] was introduced in 1976, is one of the most common data structures used in representing a graph.
  - BackLinks Compression (BLC) [2] was introduced in 2009, is a modified web-graph compression method developed by Boldi and Vigna.
  - Compressed Binary Tree (CBT) [3] was introduced in 2017, is a state-of-the-art structure, which eliminated the need for any intermediate structure to compress the graph.
- This paper introduces CSR for time-evolving graphs, also the combinations of CSR and CBT.

[1] R. A. Snay, "Reducing the profile of sparse symmetric matrices," Bulletin Ge ode sique, vol. 50, no. 4, pp. 341–352, 1976.

[2] F. Chierichetti, R. Kumar, S. Lattanzi, M. Mitzenmacher, A. Panconesi, and P. Raghavan, "On compressing social networks," in *Proceedings of the 15th ACM SIGKDD International Conference on Knowledge Dis- covery and Data Mining*, KDD '09, (New York, NY, USA), p. 219–228, Association for Computing Machinery, 2009.

[3] M. Nelson, S. Radhakrishnan, A. Chatterjee, and C. Sekharan, "Queryable Compression on Streaming Social Networks," in Big Data (Big Data), 2017 IEEE International Conference on, IEEE BigData '17, IEEE Computer Society, 2017.



#### Background – Matrix Based Compression

- Matrix based compression for time-evolving graphs,
  - In 2016, Caro et al. developed  $ck^d tree$ , by adapting the concepts of quadtree compression.
  - Evelog is a compressed adjacency log structure, based on log of events strategy.
  - Other legacy time-evolving compressions are CAS, CET, and TGCSA.
- For this paper, only  $ck^d tree$  the matrix-based compression is considered for comparison purposes, as  $ck^d tree$  has shown better results than all the other matrix-based compression.



#### Compressed Sparse Row (CSR)

- Here is an example of a Compressed Sparse Row (CSR) structure for a static undirected graph.
- In this structure it is sufficient to store only the upper triangular part of the matrix shown in green.

	0	1	2	3	4	5	6	7	8	9
0	0	1	1	1	1	1	0	0	1	0
1	1	0	0	0	1	1	1	0	0	0
2	1	0	0	0	0	0	1	1	0	0
3	1	0	0	0	0	0	0	0	0	1
4	1	1	1	0	0	0	0	0	0	1
5	1	1	1	0	0	0	0	0	0	0
6	0	1	0	0	0	0	0	0	0	0
7	0	0	0	0	0	0	0	0	0	0
8	1	0	0	0	0	0	0	0	0	0
9	0	0	0	1	1	0	0	0	0	0

Degree Array

0	1	2	3	4	5	6	7	8	9
6	3	2	1	1	0	0	0	0	0

Adjacency Array





#### Compressed Binary Tree (CBT)

- An example of a Compressed Binary Tree (CBT) structure for a static undirected graph.
- In this structure it is sufficient to store only the upper triangular part of the matrix shown in green.







# Time-evolving graph

• A time-evolving graph can be represented as several snapshots of a static graph.



- Therefore, each node in the graph has to store two sets of information,
  - One is the neighborhood array, and
  - Second, a time array to keep track of changes between each snapshot.
- Changes at each snapshot can either be,
  - Addition of an edge, or
  - Deletion of an edge, or
  - No changes at all.



#### Representation of Time-Evolving Graphs

- For example,
  - Consider node 0, which is the first row in every snapshot.
  - The time-frame array is as follows:





- Since the time instants are finite, for most of the examples, the range of the time-frames does not exceed more than 10,000.
- Hence, the time array is small compared to the entire graph.



# Compression Technique

- In this paper, the time-array and the neighborhood-array are compressed in one of two methods.
  - Compressed Sparse Row (CSR), or
  - Compressed Binary Tree (CBT).
- Either of the two methods can be chose depending up on the size and the nature of the array.
- These arrays are encoded as a series of bits.
- For the CSR, this encoding method is referred to as bit-packing.



## Bit-Packing Algorithm

- Each number in the array is represented as bits.
- Number of bits required to store each number depends on the largest number in the array.
- Then the bits are stored as an unsigned bit array.

Unsigned int	1	3	5	10	16	26	
Unsigned bit	00001 00011 00101 01010 100000 11010 00						

• A memory location can hold up-to 32 bits anymore than 32, the number is split to fit the current memory location and rest of the bits are stored in the following memory location.

Unsigned int	1	3	5	10	16	26	30		
Unsigned bit	00001 00011 00101 01010 10000 11010 11								
	110 00000 00000 00000 00000 00000 0000								



#### Input Structure

- By the definition of the time-evolving graph, each link/edge is represented as a 4-tuple (u, v, t<sub>i</sub>, t<sub>j</sub>).
- But the input can be reduced to a 3-tuple (u, v, t).
- Every odd occurrence of the triples indicates the link is active, and every even occurrence indicates otherwise.
- The input to the graph (edgelist) is first sorted with respect to the time-frames.
- For each time-frame, the edges are sorted with respect to the node number.



#### The Combinations

- Because of the two independent arrays for each node, the arrays can be stored in either be stored using CSR or CBT.
- The paper evaluates all four combinations to store the graph.
- Novel CSR CSR: Both the time-array and the neighborhood arrays are stored using CSR.
- Novel CSR CBT: The time-array is stored as CSR, and the neighborhood array is stored as CBT.
- Novel CBT CSR: The time-array is stored as CBT, and the neighborhood array is stored as CSR.
- CBT CBT: Both the time-array and the neighborhood arrays are stored using CBT.



#### Graphical Representation of CSR-CSR





CSR as unsinged char:

G@T<sub>1</sub> 000100101 100010100 110000101010001

G@T<sub>2</sub> 100010001 100100010 0010111000010101000

G@T<sub>3</sub> 0000100011000001 1111 1100111011001



# Querying operations

- Two querying operations are performed on all the algorithm on 1000 random chosen vertices.
- Given a node, and
  - A time  $t_i$ , a neighborhood query fetches all the neighbors of that node at time  $t_i$ .
  - An interval  $t_i$  to  $t_j$  (inclusive), a neighborhood query fetches all the neighbors of that node between time  $t_i$  and time  $t_j$ .
- Given nodes (u, v), and
  - A time  $t_i$ , the edge existence method returns if the edge is active/inactive at time  $t_i$ .
  - An interval  $t_i$  to  $t_j$  (inclusive), the edge existence method returns if the edge is active/inactive between time  $t_i$  and time  $t_j$ .



# Compression Result

Space required by the compression algorithms to store each graph in megabytes 10000 Size in logarithmic scale 1000 100 10 CSR - CSR CSR-CBT CBT - CSR CBT-CBT ck^d-tree .txt 271.6 CommNet 34 16 16 15.9 30 PowerLaw 546.9 80 80 70 73.8 128 860 107 73.8 82 73.8 89 ■ Flickr-Days Wiki-edits 5836.8 1433.6 1372.16 1536 2048 2048 ■ Yahoo-Netflow 19456 4403.2 2969.6 4300.8 3061.76 2560



Time required to compress each graph





#### **Experimental Results**

Graphs	Nodes	Contacts	<b>Time Frames</b>	neigh CSR Ti	neigh CSR Ti Tj	Edge exists CSR Ti	Edge exists CSR Ti -Tj
CommNet	10000	19061571	10001	$0.78\pm0.005$	$0.93\pm0.049$	$0.78 \pm 0.001$	$0.82\pm0.002$
PowerLaw	1000000	32280816	1001	$2.07\pm0.006$	$2.10\pm0.012$	$2.06 \pm 0.011$	$2.08\pm0.06$
Flickr-Days	2585570	33140018	135	$1.31 \pm 0.02$	$2.08\pm0.31$	$1.31 \pm 0.013$	$2.21 \pm 0.2$
Wiki-edits	21504191	266769613	134075025	$0.40\pm0.007$	$0.403 \pm 0.001$	$0.39 \pm 0.08$	$0.39\pm0.008$
Yahoo Netflow	32904819	1123508740	58735	$2.19 \pm 0.45$	$1.51\pm0.06$	$1.38 \pm 0.014$	$1.52 \pm 0.042$
Graphs	Nodes	Contacts	<b>Time Frames</b>	neigh_TCBT_Ti	neigh_TCBT_Ti_Tj	Edge exists_CBT_Ti	Edge exists_CBT_Ti-Tj
CommNet	10000	19061571	10001	$1.33 \pm 1.44$	$1.43\pm0.47$	$0.39 \pm 0.66$	$0.39\pm0.55$
PowerLaw	1000000	32280816	1001	$2.99\pm0.55$	$5.70 \pm 1.05$	$0.64 \pm 0.12$	$1.6\pm0.03$
Flickr-Days	2585570	33140018	135	$11.29 \pm 8.52$	$38.49 \pm 8.34$	$4.23 \pm 2.81$	$5.44 \pm 10.11$
Wiki-edits	21504191	266769613	134075025	1.24, 1.911	1.42, 2.12	1.15, 0.18	1.15, 0.19
Yahoo Netflow	32904819	1123508740	58735	43.13, 0.263	51.21, 4.67	30.32, 2.36	31.2, 5.37
				-			
Graphs	Nodes	Contacts	<b>Time Frames</b>	neigh CKD Ti	neigh CKD Ti Tj	Edge exists CKD Ti	Edge exists CKD Ti -Tj
CommNet	10000	19061571	10001	$48.89 \pm 11.56$	$64.46\pm0.43$	$49.6 \pm 3.4$	$49.7\pm0.24$
PowerLaw	1000000	32280816	1001	$374.23 \pm 50.72$	$374.64 \pm 50.66$	$216.0 \pm 5.3$	$226.13 \pm 14.38$
Flickr-Days	2585570	33140018	135	$35.34 \pm 10.39$	$45.22 \pm 5.78$	$35.2 \pm 1.2$	$37.2 \pm 2.3$
Wiki-edits	21504191	266769613	134075025	$3.0 \pm 3.0$	$4.39 \pm 0.72$	$2.62 \pm 1.7$	$2.98 \pm 0.25$
Yahoo Netflow	32904819	1123508740	58735	$231.9 \pm 82.1$	$254 \pm 92.06$	$211.8 \pm 89.0$	$212.32 \pm 71$

All the experiments were run on an Intel(R) Xeon(R) CPU E5520 @ 2.27GHz (16 Cores) with 64 GB of RAM, and the programs are written in GNU C/C++.



#### Conclusion

- Valuable insights can be gained from the analysis of time-evolving graphs.
- Our techniques show a significant reduction in memory requirement.
- All algorithms are tested on real-world datasets and show significant improvements over existing techniques.
- Our future work would focus on exploiting the parallelism in improving the compression techniques' timings in a broader domain of graphs..





# Thank you.

Question?