The Seventh International Conference on Advances in Databases, Knowledge, and Data Applications
May 24-29, 2015 - Rom, Italy

# Overview of Regular Path Queries in Graphs 

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## Outlook

- Introduction
- Application Fields
- Types of Graphs
- Regular Path Queries (RPQ)
- Extensions of RPQ


## Introduction

- Query languages for graph databases ...
- are navigational/recursive
- traverse the labeled edges/nodes
- query the topology of the data, but not the data itself
- Basic building blocks for this languages are often regular path queries (RPQ)
- RPQ are used in many graph query languages, so i.e. in
- G
- GraphLog
- Cypher
- XPath
- SPARQL 1.1


## Application Fields for Graphs and Regular Graph Queries

- Knowledge representation (RDF/s, OWL, Ontologies like Yago, Taxonomies)
- connection between entities, instance and subclass relationships
- Transportation networks (airline, train, bus, streets)
- Reachability, shortest path, critical path
- geographical information
- biological applications
- bibliographic citation analysis
- social networks
- program analysis
- reachability of code, variables used before defined, deadlocks


## Query Types

- Graph Pattern Matching (subgraph matching)
- Path Finding (finding nodes connected by graphs)
- Extraction of edge label variables
- Aggregation
- ...


## Types of Graphs

- Undirected
- Cyclic


Distances from http://www.luftlinie.org

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## Types of Graphs



## Property Graph

- directed
- cyclic
- Nodes and relations have properties



## Yet another Graph: Type ?



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## Definition of Database Graph

$\mathrm{G}=(\mathrm{N}, \mathrm{E})$
Directed, labeled graph with:

- $N$ : set of nodes

$\mathrm{G}=(\mathrm{N}, \mathrm{E})$
Directed, labeled graph with:
- N : set of nodes, representing entities from the real world


## Country

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Germany
```

Karlsruhe Rom Koper Tokyo

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## Definition of Database Graph [MW89]

$\mathrm{G}=(\mathrm{N}, \mathrm{E})$
Directed, labeled graph with:

- N : set of nodes,, representing entities from the real world
- $E$ : set of directed edges


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## Definition of a Database Graph for RQP

$G=(N, E)$
Directed, labeled graph with:

- N: set of nodes
- E: set of directed edges
- S: finite set of symbol for labeling of edges (vocabulary)



## Regular Path Queries (RPQ)

- RPQ have the form: RQP( $\mathrm{x}, \mathrm{y}$ ) := (x, R, y) where $R$ is a regular expression over the vocabulary of edge labels
- Construction of regular expressions: $R::=s|R . R| R|R| R^{*}|R ?|(R) \quad / / ~ s$ element from $S$
- Examples:
- Ancestors: isChildOf+
- Cousins
isChildOf . isMarriedWith . isChildOf . hasChild . hasChild

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## RPQ Example

$$
\mathrm{R}=\mathrm{a}+(\mathrm{d} \mid \mathrm{b}) \mathrm{ab}
$$



RPQ Example


$$
\begin{aligned}
& \boldsymbol{R}=\mathrm{a}+(\mathrm{d} \mid \mathrm{b}) \mathrm{ab} \\
& \text { adab --> }(2,10)
\end{aligned}
$$

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RPQ Example in informacijske tehnologije

$$
\begin{aligned}
& R=a+(d \mid b) a b \\
& \text { adab --> }(2,10) \\
& \text { aadab --> }(1,10)
\end{aligned}
$$

RPQ Example


$$
\begin{aligned}
& R=a+(d \mid b) a b \\
& \text { adab --> }(2,10) \\
& \text { aadab --> }(1,10) \\
& \text { abab --> }(3,10)
\end{aligned}
$$

## Algorithms for Answering RPQ

- Mapping to finite automaton [Mendelzon, Wood, 1995]
- Construct finite nondeterministic automata from query with start state $\mathrm{s}_{0}$ und final state $\mathrm{s}_{\mathrm{f}}$
- Consider G as NFA with start state x and final state y
- Form product automaton
- Determine if there is a path form $\left(\mathrm{s}_{0}, \mathrm{x}\right)$ to $\left(\mathrm{s}_{\mathrm{f}}, \mathrm{y}\right)$
- Search for rare labels and start BFS [Koschmieder, 2012]
- Look for mandatory rare labels in the query (concering the graph)
- Use the nodes from the rare edge labels as starting points for a two-way search between endpoints and startpoints of the rare edges


## Two-way Regular Path Queries (2RQP)

- 2RPQ extend the vocabulary of RPQ by the „inverse" of each relationship symbol.
- Foreach symbol s in S: there exist a symbol s-
- Example:


$$
\mathrm{R}=\mathrm{a}+\mathrm{d}^{-} \mathrm{ab}^{-}
$$

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## 2RPQ Example

$$
R=a+d^{-} a b^{-}
$$



Results: $(2,1)$

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## 2RPQ Example

$$
R=a+d^{-} a^{-}
$$



Results: $(2,1)$

## Conjunctive Regular Path Queries (CRPQ)

- $\operatorname{CRPQ}\left(z_{1}, \ldots, z_{n}\right)=\operatorname{AND}_{(1<=i<=m)}\left(x_{i}, R_{i}, y_{i}\right) \quad / /$ each $z_{i}$ is a $x_{i}, y_{i}$
- Examples:
- Which couples are married by a pontifex?

CRPQ $(x, y, p):=(x$, isMarriedWith, $y)$ AND
( $x$, isMarriedBy, $p$ ) AND
( y , isMarriedBy, p ) AND
( $p$, isa, 'Pontifex')

- Related at mostly over '5 edges' (navigating the family tree)

CRPQ $(x, y):=(x$, isChildOf\{5\}, z) AND
( $y$, isChildOf\{5\}, $z$ )

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## CRPQ Example

( $\mathrm{x}, \mathrm{a}+\mathrm{y}$ ) AND ( $\mathrm{x}, \mathrm{e}+\mathrm{y}, \mathrm{y})$


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## CRPQ Example

$$
(x, a+, y) \text { AND }(x, e+, y)
$$



Result: (2, 8) in informacijske tehnologije

## Extended Conjunctive Regular Path Queries (ECRPQ)

- CRPQ extended by
- allow free path variables in the query
- checking relations on sets of paths
- Example 1: Return all paths between $x$ and $y$, which have a concrete node e (id:123) in between:
- ECRPQ $\left(p_{1}, p_{2}\right)=\left(x, R_{1}, e\right)$ AND $\left(e, R_{2}, y\right)$ AND (e, hasID, 123)
- Example 2: Path pattern match - Find all node connected by paths of the form $a^{n} b^{n} c^{n}$ :
- $\left.\operatorname{ECRPQ}(x, y)=\left(x, p v_{1}, z_{1}\right),\left(z_{1}, p v_{2}, z_{2}\right),\left(z_{2}, \operatorname{pv}_{3}\right) y\right)$,



## Aggregation (AggCRPQ)

- CRPQ + Aggregation functions, i.e. for calculationg the distance between nodes
- Examples:
- How many biological children does the husband of Carolyn has?:

AggCRPQ $(x, \operatorname{count}(y))=(x$, isMarriedWith, 'Carolin'), $(x$, isParentOf, $y)$

- Shortest path between $x$ and $y$ (with intermediate node $z$ )
$\operatorname{AggCRPQ}\left(x, y, \min \left(\operatorname{len}\left(p_{1}\right)+\operatorname{len}\left(p_{2}\right)\right)\right)=\left(x, p_{1}, z\right),\left(z, p_{2}, y\right)$


## Summary \& Outlook

- RPQ and its extensions are partly/complete realized in a number of graph query languages
- Different extensions of RPQ provide additonal power of expressiveness
- In most implementations of graph query languages RPQ are combined with additional data query functionalities
- Complexity and Containment is actual research field


## Literature

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- SPARQL Property Paths: http://www.w3.org/TR/sparql11-property-paths/ in informacijske tehnologije


## SPARQL 1.1 path language

| Syntax Form | Matches |
| :---: | :---: |
| uri | A URI or a prefixed name. A path of length one. |
| $\wedge$ elt | Inverse path (object to subject). |
| (elt) | A group path elt, brackets control precedence. |
| elt1 / elt2 | A sequence path of elt1, followed by elt2 |
| elt1 ^ elt2 | Shorthand for elt1 / ^elt2, that is elt1 followed by the inverse of eltz |
| elt1 \| elt2 | A alternative path of elt1, or elt2 (all possibilities are tried). |
| elt* | A path of zero or more occurrences of elt. |
| elt+ | A path of one or more occurrences of elt. |
| elt? | A path of zero or one elt. |
| elt ( $n, m$ ) | A path between n and m occurrences of elt. |
| elt ( $n$ ) | Exactly $n_{n}$ occurrences of elt. A fixed length path. |
| elt $\{n$, \} | $n$ or more occurrences of elt. |
| elt $\{, \mathrm{n}$ ) | Between 0 and $n$ occurrences of elt. |

