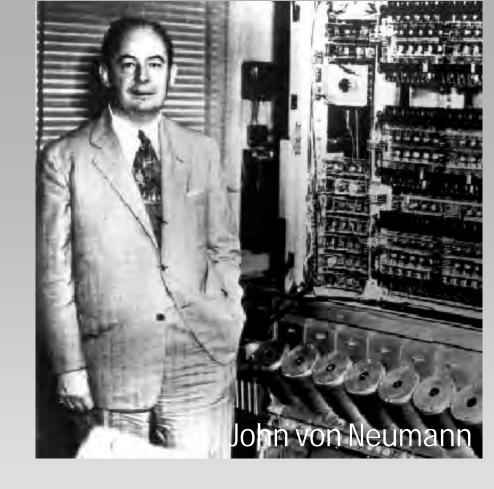
Linguistic Geometry: Adversarial Reasoning for Real Life Problems

Boris Stilman University of Colorado Denver, USA STILMAN Advanced Strategies, USA

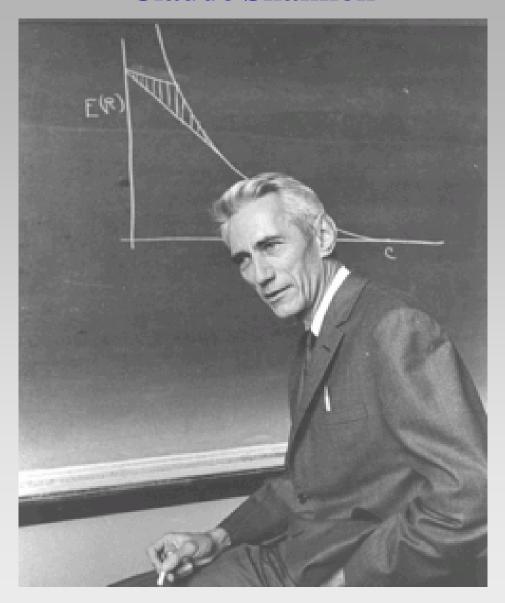
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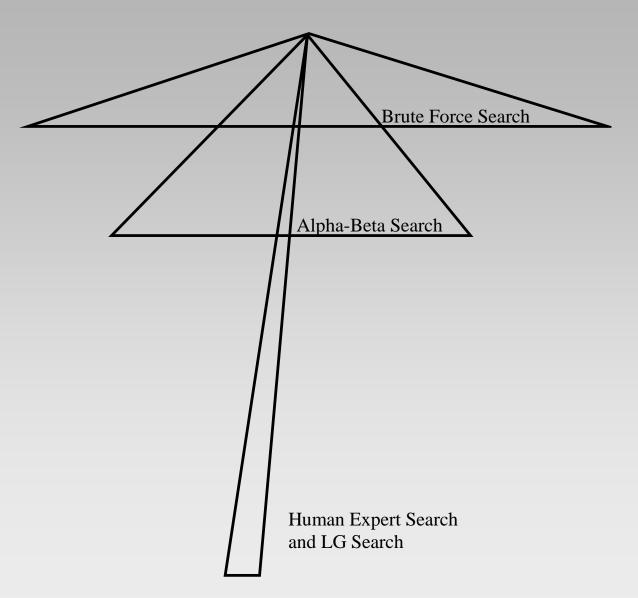


What is Linguistic Geometry?

Claude Shannon



A Comparison of Searches for the same processing time



1997

Rematch
Garry Kasparov
vs
Deep Blue

Facts Garry Kasparov IBM's Deep Blue

Height: 5'10" 6'5"

Weight: 176 lbs. 1.4 tons Age: 34 years 4 years

Birthplace: Azerbaijan Yorktown, NY

Processors: 100B Neurons 32 P2SC Processors

Moves/Second: 2 200 million

Power Source: electrical/chemical electrical

Next Career: champion none

Chess World Champions Mikhail Botvinnik



1911 - 1995 6th World Champion 1948 - 1957, 1958 - 1960, 1961 - 1963



Mikhail Botvinnik
"To Achieve the Goal"





Moscow State University АНАДЕМИЯ НАУН СССР ИЛУЧНЫЙ СОВЕТ ПО КОМПЛЕКСНОЙ ПРОБЛЕМЕ "КИБЕРНЕТИНА"

м. м. ботвинник

БЛОК-СХЕМА АЛГОРИТМА ИГРЫ В ШАХМАТЫ



MOCHBA 1972

M. Botvinnik

A Flow Chart of The Algorithm for Playing Chess Depreciation Schedules • Income Taxes • TRS-80 Electric Pencil

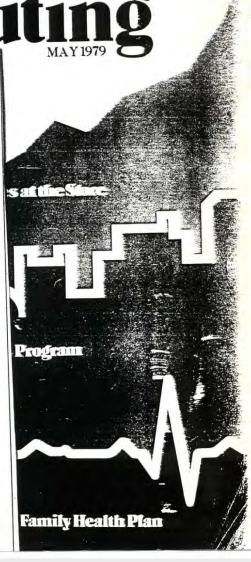
Personal Computing There have been a few challenges to MAY 1979

... There have been a few challenges to the "Botvinnik Computer Problem." Typical of these challenges is the letter from Thomas Chassereau, 1345 53rd Ave., Oakland, CA 94601. "I think your magazine is of great interest," writes Tom. "And I enjoy reading it, not only for enjoyment but as a supplement to my college education with a major in Data Processing and a minor in

Math. I have been playing chess for over 20 years and I am astounded by the solution given to the Botvinnik End Game, as solved by his PIONEER program!" (See the diagram for original position.)



Original position of Botvinnik's endgame problem with White to win as solved by the PIONEER Computer program. (1. Pg5-g6; Kf5-f6;1. Pg6g7; Kf6-f7; 3. Pg7-g8=Q and White

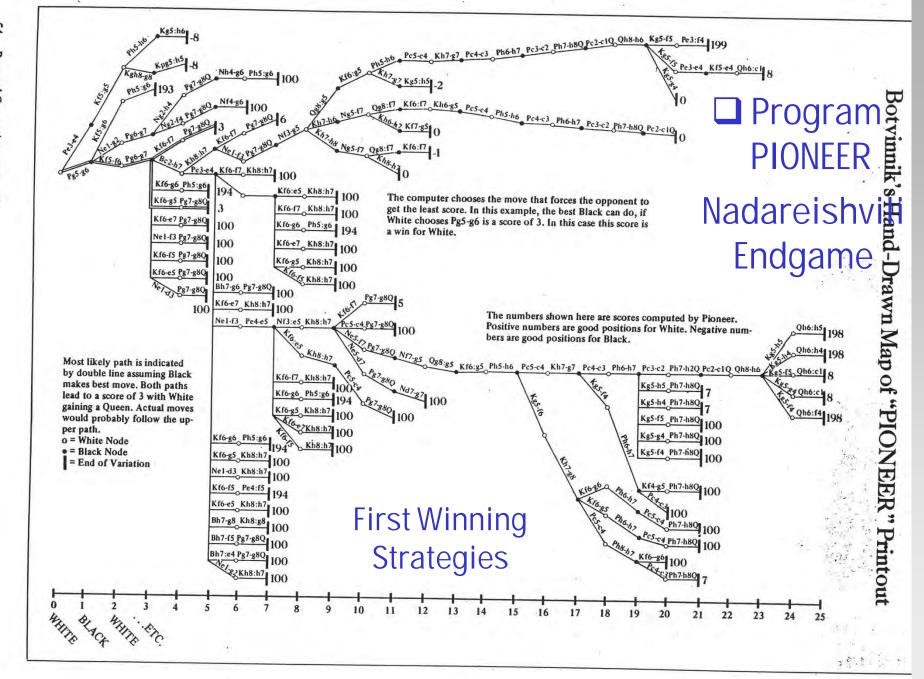


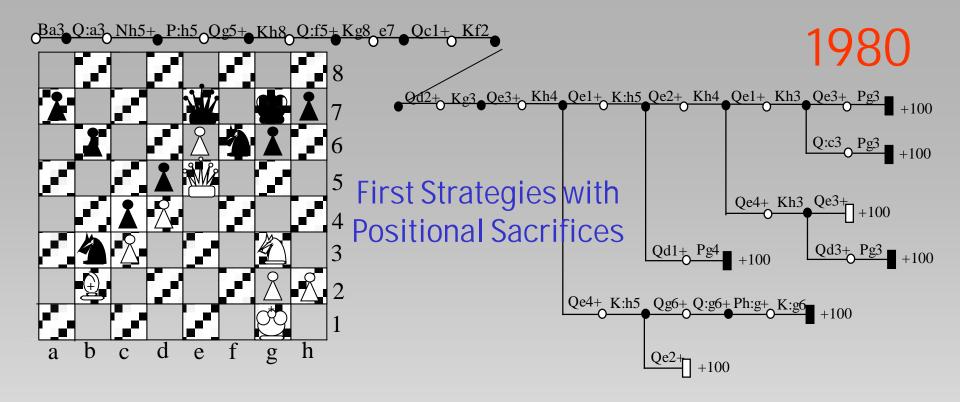
1977

☐ Program PIONEER

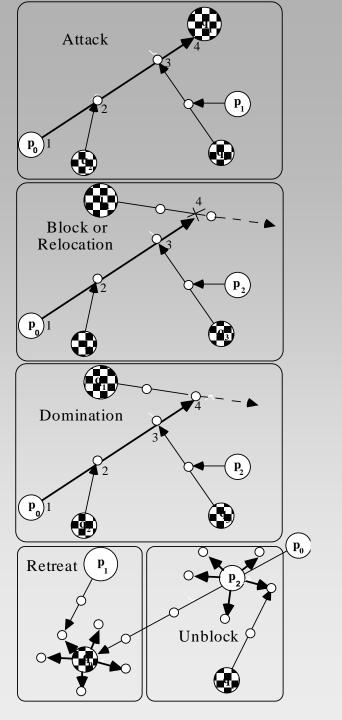
Nadareishvili Endgame







Moscow, USSR Program PIONEER



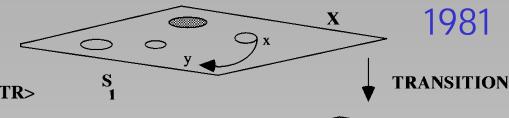
Zones

Types of Zones

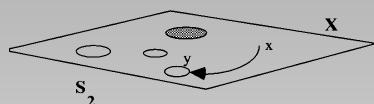
Abstract Board Games



$$<\mathbf{X},\mathbf{P},\mathbf{R}_{p},\left\{ \mathbf{ON}\right\} ,\mathbf{v},\mathbf{S}_{i},\mathbf{S}_{t},\mathbf{TR}>$$



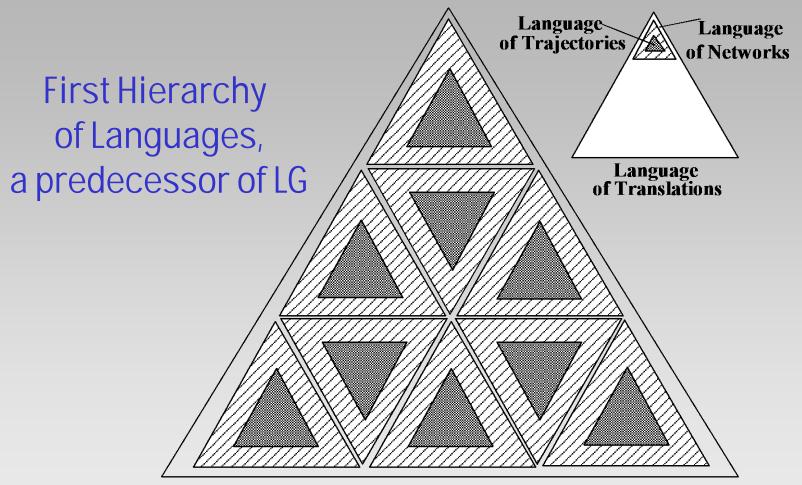
$\mathbf{X} = \{\mathbf{x_i}\}$	is a finite set of <i>points</i> ;
$P = \{p_i\}$	is a finite set of <i>elements</i> ; $P=P_1 \cup P_2, P_1 \cap P_2=\{\};$
$\mathbf{R}_{\mathbf{p}}(\mathbf{x}, \mathbf{y})$	is a family of binary relations of <i>reachability</i> in X $(x \in X, y \in Y, p \in P)$; y is <i>reachable</i> from x for p;



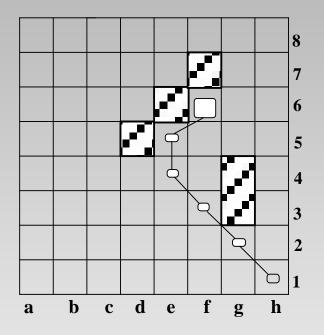
- ON(p) = x is a partial function of *placement* of elements P into X;
- $\mathbf{v} > 0$ is a real function, $v(p_i)$ are the *values* of elements;
- S_i is a set of *initial* states of the system, a certain set of formulas $\{ON(p_i)=x_i\}$;
- S_t is a set *target* states of the system (as S_i);
- TR is a set of operators TRANSITION(p, x, y) for transition of the system from one state to another described as follows

precondition:
$$(ON(p) = x) \land Rp(x, y)$$

delete: $ON(p) = x$
add: $ON(p) = y$

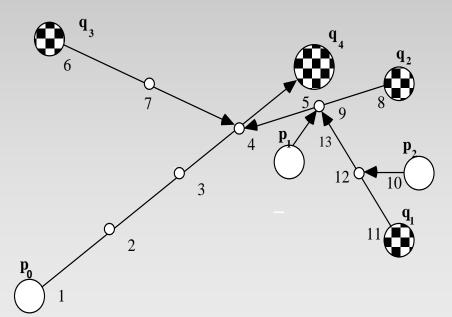


a(f6)a(e5)a(e4)a(f3)a(g2)a(h1)



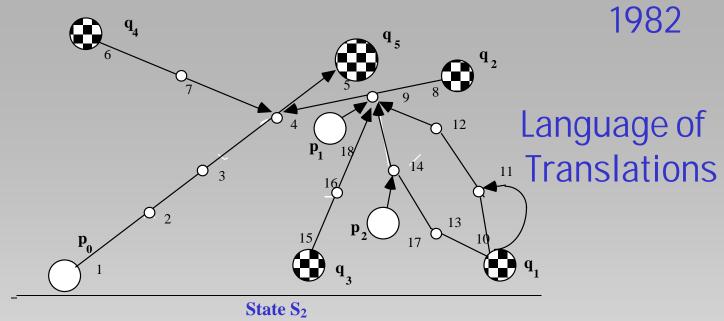
Language of Trajectories

$$Z = t(p_0, a(1)a(2)a(3)a(4)a(5), 5)t(q_3, a(6)a(7)a(4), 4)$$
$$t(q_2, a(8)a(9)a(4), 4)t(p_1, a(13)a(9), 3)$$
$$t(q_1, a(11)a(12)a(9), 3)t(p_2, a(10)a(12), 2)$$

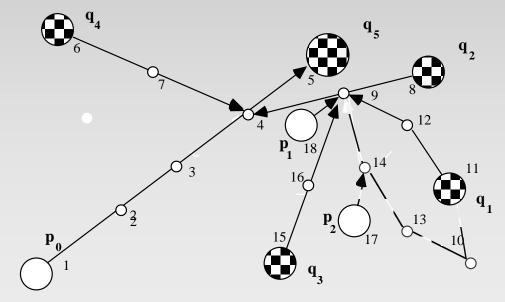


Language of Networks (Zones)

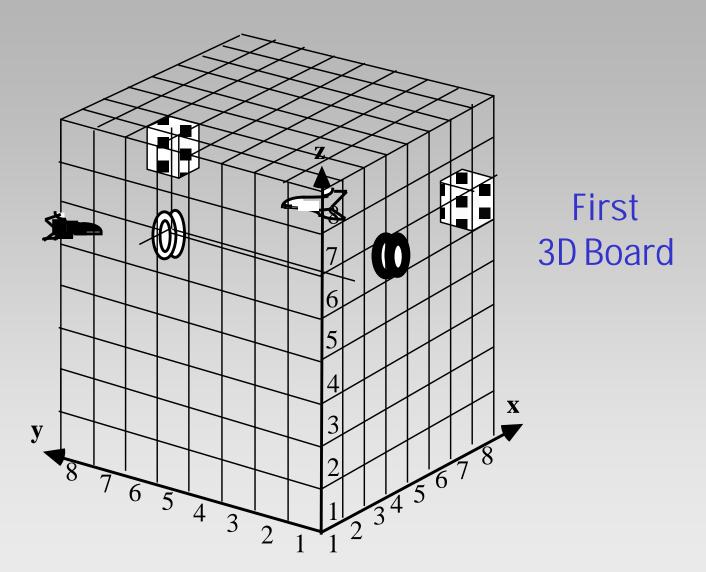
State S_1

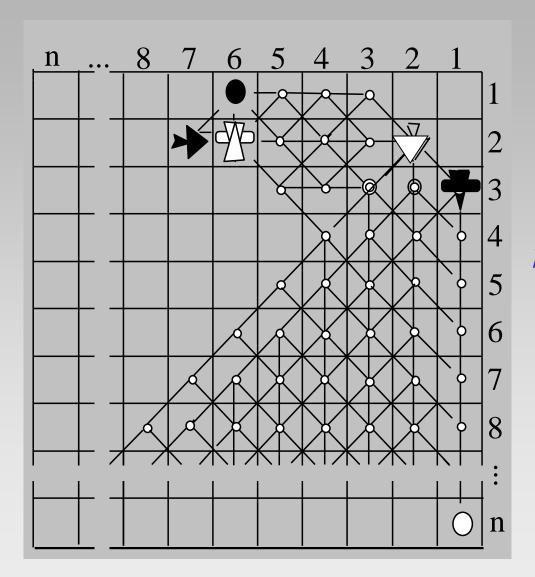




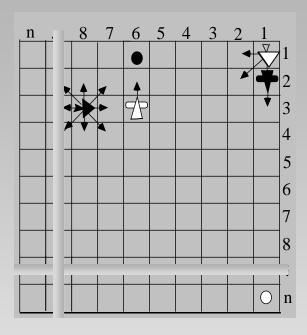




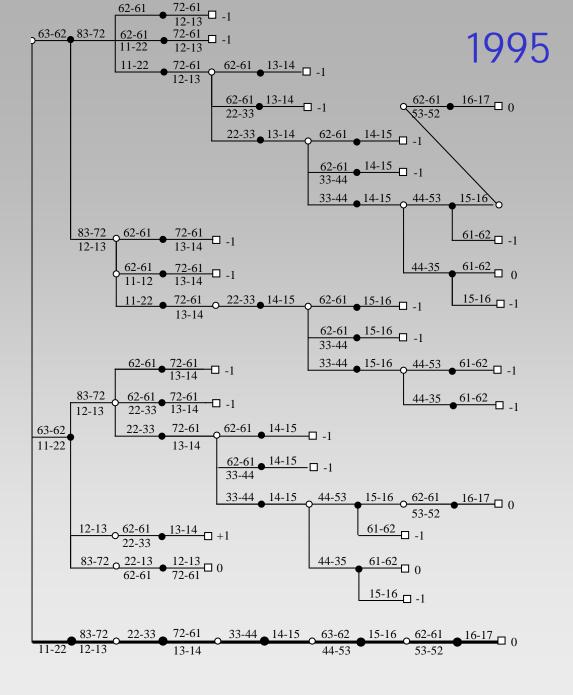




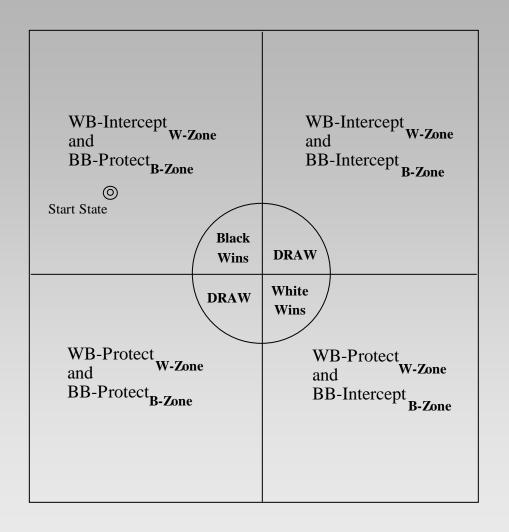
First $n \times n$ Board



First Concurrent Strategies

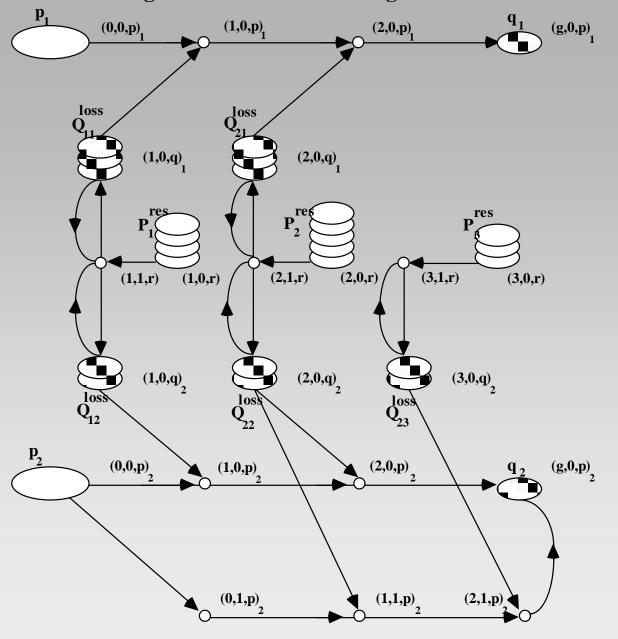


Experiments with State Space Chart: NO-Search Approach



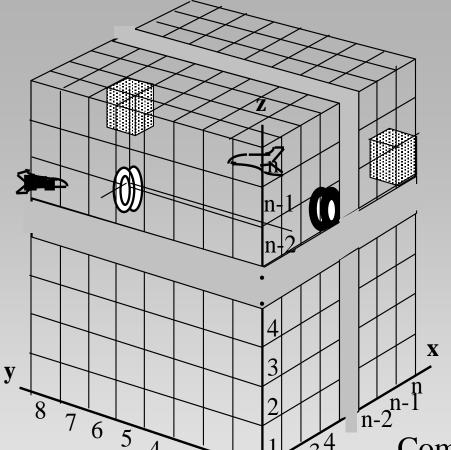
First Proof of Strategies Optimality

Planning of Maintenance Planning of Power Units



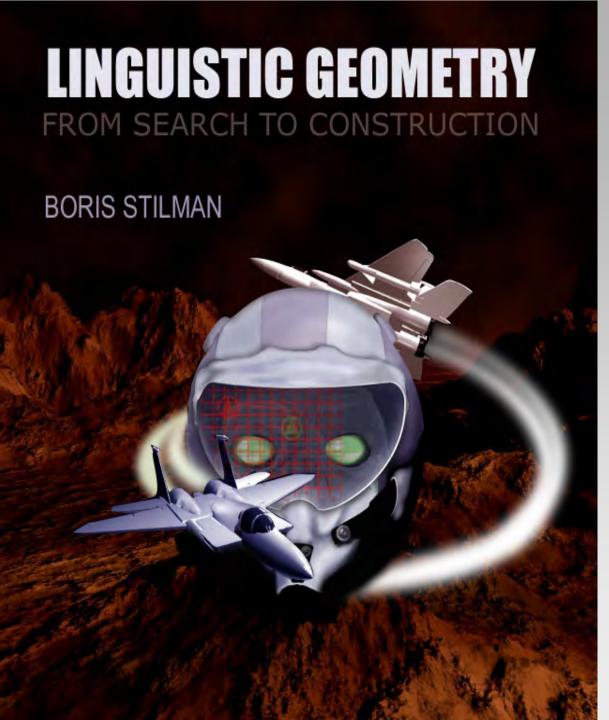
1997

First Artificial Enemy in LG



First Proof of Polynomial Run Time

Computational complexity of the specific classes of abstract board games is **polynomial** with respect to the length of the input.

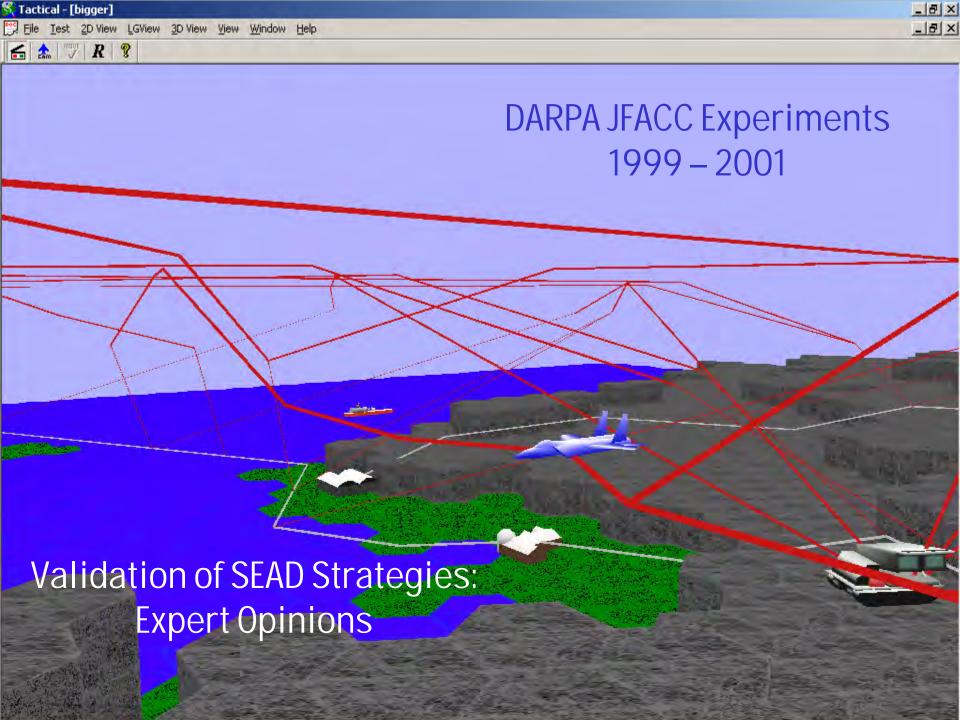


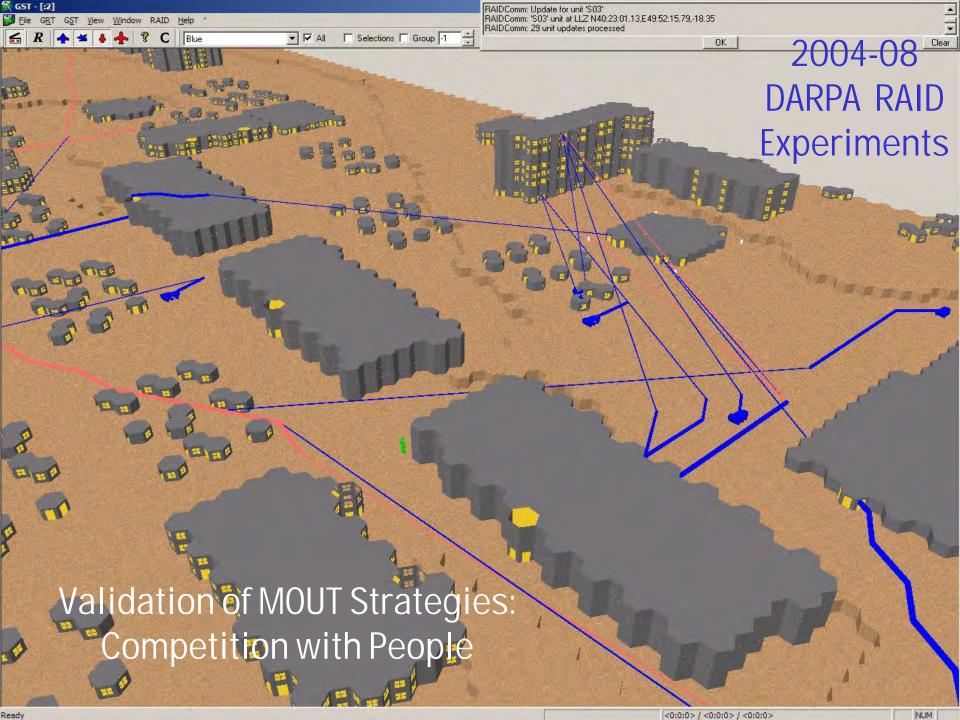
Manuscript finished in 1999 Published in 2000

- Linguistic Geometry (LG) is a new type of game theory;
- LG replaces search by construction, making the games computationally tractable.

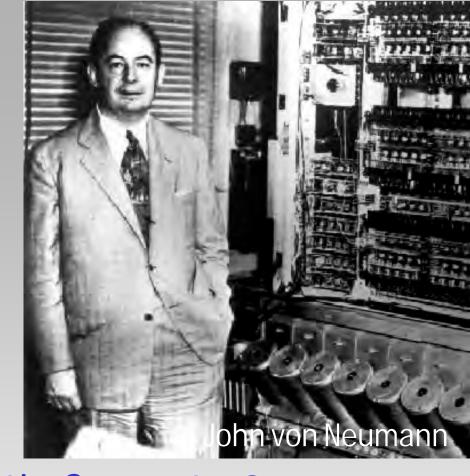


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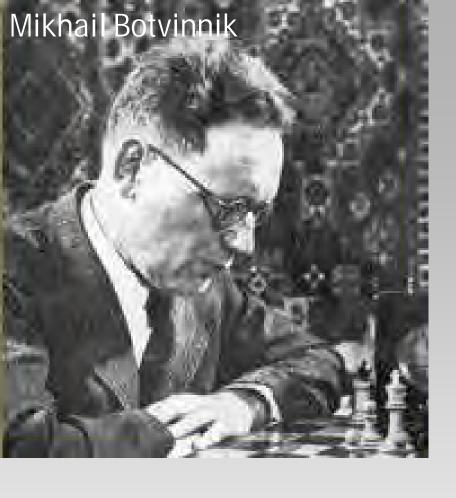


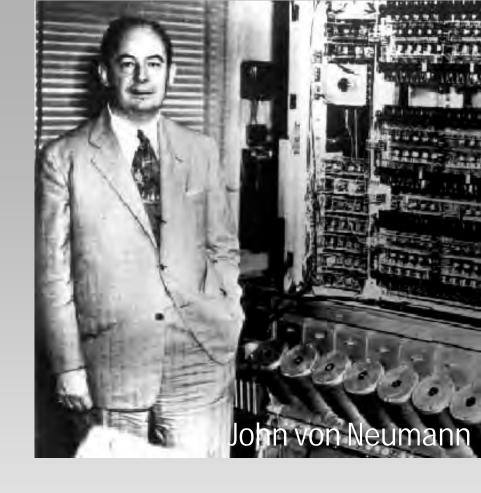




What is Linguistic Geometry?

- Chess Masters' Problem Solving
- Search Reduction Techniques
- Subclass of Gaming Problems of Polynomial Complexity
- No-Search Paradigm for Decision Making: "From Search to Construction"
- OR maybe, it is something else . . .





It appears that LG is a model of human thinking about conflict resolution, a warfighting model at the level of superintelligence . . .